


Measurement-Based Quantum Computation in Realistic Spin-1 Chains

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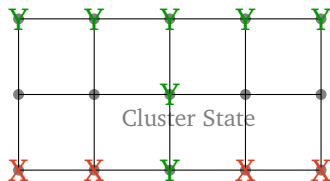
Perimeter Institute

Measurement-Based Quantum Computation

Single-site mmts on an entangled state

instead of

unitary gates on an unentangled state



- New avenue in QC research — which states are useful?
- Practical implications — can we engineer these states?
- Connections to condensed matter — maybe ground states of realistic Hamiltonians are useful?

Explore the AKLT model

Outline:

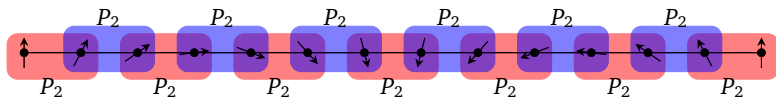
- 1 The AKLT model and how to perform single-qubit logic
Brief look at physical implementation
- 2 How to characterize computational ability of imperfect states
Numerical results: Computational ability is robust!
- 3 Improving performance: gate buffering and AKLT “renormalization”
- 4 Summary and Outlook



Spin-1 Chain

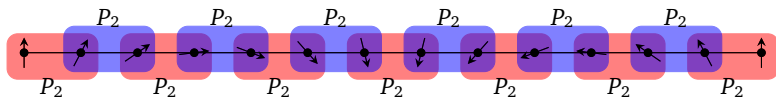
Affleck-Kennedy-Lieb-Tasaki: Consider the nearest-neighbor Hamiltonian

$$H_{\text{AKLT}} = \sum_{j=1}^n (\vec{S}_j \cdot \vec{S}_{j+1}) + \frac{1}{3} (\vec{S}_j \cdot \vec{S}_{j+1})^2 \simeq \sum_{j=1}^n (P_2)_{jj+1}$$



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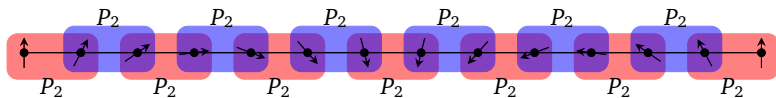
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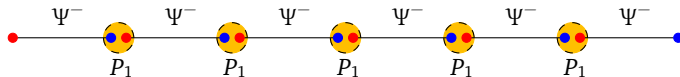
- ☞ Ground state is unique periodic BCs; 4fold degenerate open BCs
or $n \rightarrow \infty$ and $n < \infty$
- ☞ Gap to first excited state (conjectured by Haldane, analytic example by AKLT)
- ☞ Ground state is a “valence bond solid” (VBS), frustration-free

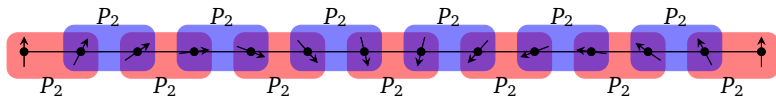


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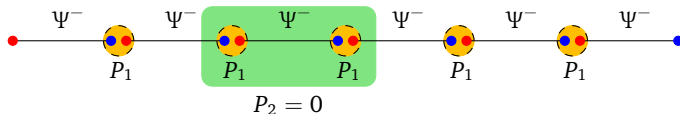




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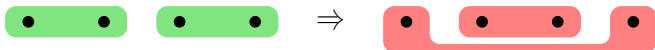
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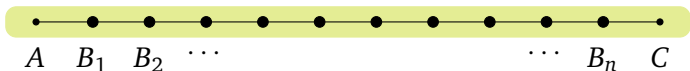


VBS is the original matrix-product state

- Easy to work out using entanglement swapping relations



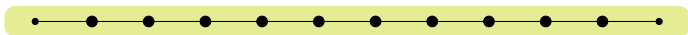
- For $|s\rangle \equiv |m_s = 0\rangle$, $s = x, y, z$ & Pauli operators σ_s (but with $\sigma_y = \sigma_x \sigma_z$):



$$|\mathcal{G}_{\text{open}}\rangle = \sum_{\{s_k\}} |s_1, s_2, \dots, s_n\rangle_B \otimes (\sigma_{s_n} \sigma_{s_{n-1}} \cdots \sigma_{s_1})_C |\Psi^-\rangle_{AC}$$

$$|\mathcal{G}_{\text{periodic}}\rangle = \sum_{\{s_k\}} \text{Tr} [\sigma_{s_n} \sigma_{s_{n-1}} \cdots \sigma_{s_1}] |s_1, s_2, \dots, s_n\rangle$$

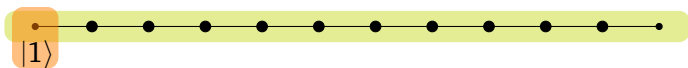
Single-qubit logic on the VBS



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☞ Chain encodes one logical qubit (think of it at C)

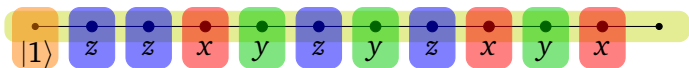
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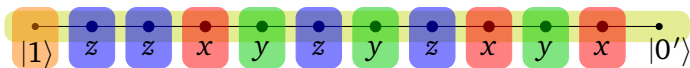
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- ☞ To do nothing, measure in the $|s\rangle$ basis
 “Byproduct” operators accumulate on C (change of basis)

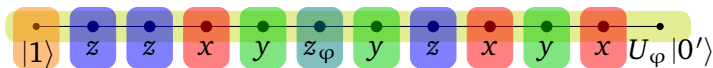
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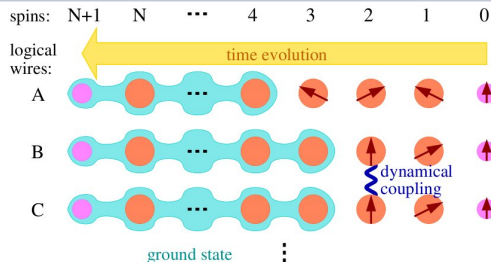
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- Chain encodes one logical qubit (think of it at C)
- Initialize: Measure $|0\rangle, |1\rangle$ on end qubit A
- To do nothing, measure in the $|s\rangle$ basis
“Byproduct” operators accumulate on C (change of basis)
- To rotate, measure in a rotated basis. Example: $R_{\hat{x}}(\varphi)$
 $\{|x_\varphi\rangle = |x\rangle, |y_\varphi\rangle = \cos \frac{\varphi}{2} |y\rangle + i \sin \frac{\varphi}{2} |z\rangle, |z_\varphi\rangle = i \sin \frac{\varphi}{2} |y\rangle + \cos \frac{\varphi}{2} |z\rangle\}$.
 Byproducts are σ_x, σ_y , and σ_z , respectively.
- Rotation is probabilistic, but heralded.

Brennen & Miyake scheme PRL 101, 010502 (2008).



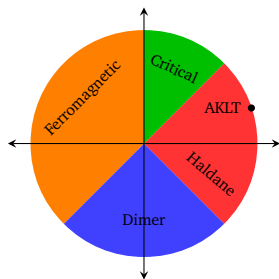
- 2D lattice of spins
- Individual chains as “quantum wires”
- Two-qubit logic (CPHASE) by dynamical coupling

Physical Implementation

Atoms or polar molecules in an optical lattice

Gap protects against thermal noise

Imperfect Hamiltonians



☞ Can only engineer H_{AKLT} to some accuracy.

☞ Consider the Hamiltonians

$$H(\theta) = \sum_{j=1}^n \cos \theta (\vec{S}_j \cdot \vec{S}_{j+1}) + \sin \theta (\vec{S}_j \cdot \vec{S}_{j+1})^2.$$

☞ System is gapped in the Haldane phase (Haldane's conjecture/fact)

☞ Computation should still be robust to noise.

☞ Using the same measurement scheme as for the AKLT state, how well can we perform single-qubit operations?

- Can in principle determine the output fidelity for any given mmt sequence and initial state. \Rightarrow But there are an exponential number!

- Work in the gate teleportation picture (Doherty & Bartlett PRL 103, 020506 (2009).)



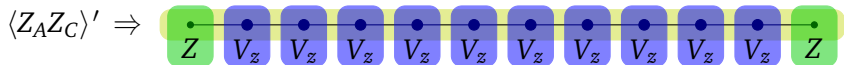
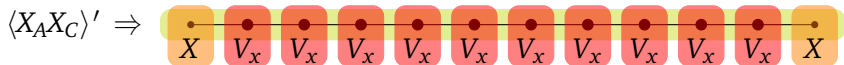
$$|\mathcal{G}\rangle \Rightarrow |s_1, s_2, \dots, s_n\rangle_B \otimes (\sigma_{s_n} \sigma_{s_{n-1}} \cdots \sigma_{s_1})_C |\Psi^-\rangle_{AC}$$

- Now only need to evaluate $\langle X_A X_C \rangle_{\vec{s}}$ and $\langle Z_A Z_C \rangle_{\vec{s}}$ for each mmt result \vec{s} , suitably corrected for byproducts, and average ($|\Psi^-\rangle$ gives -1 for both)

- $\langle X_A X_C \rangle' = \sum_{\vec{s}} \langle X_A X_C \rangle_{\vec{s}}$ with $\langle X_A X_C \rangle_{\vec{s}} = \langle \mathcal{G} | X_A P_{\vec{s},B} U_{\vec{s},C}^\dagger X_C U_{\vec{s},C} | \mathcal{G} \rangle$.
 $P_{\vec{s}}$ is the measurement projector, and $U_{\vec{s}}$ is the correction operator.

- Want to evaluate $\langle X_A X_C \rangle_{\vec{s}} = \langle \mathcal{G} | X_A P_{\vec{s},B} U_{\vec{s},C}^\dagger X_C U_{\vec{s},C} | \mathcal{G} \rangle$
- $U_{\vec{s}}$ is a Pauli operator, just $X^{\#(x \text{ or } y)} Z^{\#(z \text{ or } y)}$.
 - $(Z^{\#(z \text{ or } y)} X^{\#(x \text{ or } y)}) X (X^{\#(x \text{ or } y)} Z^{\#(z \text{ or } y)}) = (-1)^{\#(z \text{ or } y)} X$
 - $(-1)^{\#(z \text{ or } y)} |s\rangle = e^{i\pi J_x} |s\rangle$.
 - $\langle X_A X_C \rangle_{\vec{s}} = \langle \mathcal{G} | X_A \vec{V}_{x,B} P_{\vec{s},B} X_C | \mathcal{G} \rangle$ for $\vec{V}_{x,B} = (\bigotimes_{k=1}^n e^{i\pi J_{x,B_k}})$.

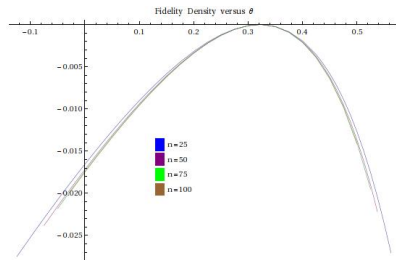
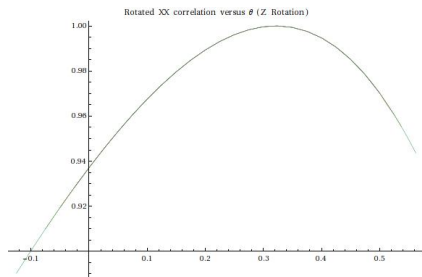
Great! Now $\sum_{\vec{s}} P_{\vec{s},B} = \mathbb{1}_B$ and $\langle X_A X_C \rangle' = \langle \mathcal{G} | X_A \vec{V}_{x,B} X_C | \mathcal{G} \rangle$



- ➡ We only need to evaluate string operators $X_A \vec{V}_{x,B} X_C$ and $Z_A \vec{V}_{z,B} Z_C$.
- ☞ Similar to string operators $S_{x,j} \vec{V}_x S_{x,k}$ showing the “hidden order” of ground states in the Haldane phase (finite at any length, unlike usual spin correlations which decay exponentially.)
- ☞ Haldane conjecture/fact (plus fixed BCs) $\Rightarrow \langle X_A X_C \rangle' = \langle Z_A Z_C \rangle' = 1$:
 - $X_A \vec{V}_{x,B} X_C$ is a global rotation by π about \hat{x} ;
 - H is rotationally-invariant and ground state unique.
- ➡ Near-AKLT ground states also have well-defined qubit encoding.
- ☞ For single qubit rotations we need $\langle X_A R(\theta) X_C R(\theta)^\dagger \rangle'$, etc. Attempt rotations only a few times, average results.

Numerical results:

- ☞ Find ground states of the open/fixed chain using DMRG (actually use MPS form and imaginary time evolution)
- ☞ Localizable entanglement = 1 throughout Haldane phase, as expected.
- ☞ Rotation fidelity drops off slowly with theta (no computational phase)



- Nearby Hamiltonian: $H(\theta_{\text{AKLT}} + \epsilon) \approx H_{\text{AKLT}} + \frac{5}{6}\epsilon \sum_{j=1}^n \text{SWAP}_{j,j+1}$
- Perhaps $|\mathcal{G}(\theta)\rangle = \alpha |\mathcal{G}_{\text{AKLT}}\rangle + \sum_{j=1}^n \beta_j \text{SWAP}_{j,j+1} |\mathcal{G}_{\text{AKLT}}\rangle$?
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 - Doing nothing is already swap invariant: only # of x 's, y 's, z 's matters.
 - $R_x(\varphi)$: Action of mmt sequence $|x\rangle |z_\varphi\rangle |x\rangle \Rightarrow X(ZR_x(\varphi))X = -ZR_x(\varphi) |z_\varphi\rangle |x\rangle |x\rangle$, $|x\rangle |x\rangle |z_\varphi\rangle \Rightarrow ZR_x(\varphi)X$. Swap invariance! (up to a sign)
 - For $|\mathcal{G}(\theta)\rangle$ a superposition of $|\mathcal{G}_{\text{AKLT}}\rangle$ and $\text{SWAP} |\mathcal{G}_{\text{AKLT}}\rangle$, there are different paths to same properly-rotated state, but ± 1 phase.

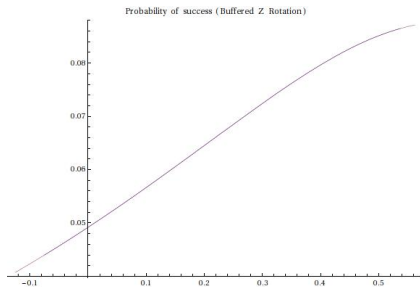
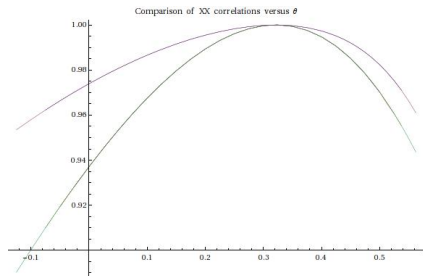
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- ⇒ Fidelity should improve, but overall probability of success will vary



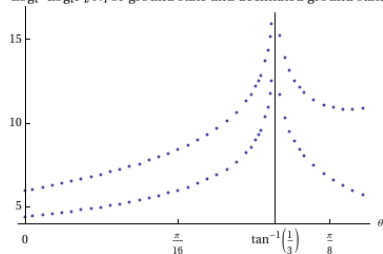
☞ Is there some simple transformation making the chain more AKLTlike?

- Should have something to do with repeated mmt outcomes
- Should respect the rotational symmetry

☞ Project neighbors onto spin-zero — $|\mathcal{G}_{\text{AKLT}}\rangle$ invariant!

- $|S = 0\rangle = \frac{1}{\sqrt{3}} (|xx\rangle - |yy\rangle + |zz\rangle)$
- $|\mathcal{G}\rangle = \sum_{\{s_k\}} |s_1, s_2, \dots, s_n\rangle_B \otimes (\sigma_{s_n} \sigma_{s_{n-1}} \cdots \sigma_{s_1})_C |\Psi^-\rangle_{AC}$

-Log[-Log[F]/N] of ground state and decimated ground state



☞ Project neighbors onto spin 0 or 1

☞ Project n-n-neighbors onto spin-zero (gate buffering)

☞ Probably many more

☞ Dimer state also a fixed-point.

Summary and Outlook

- ☞ Can efficiently characterize computational ability of spin-1 chains.
 - ☞ Near-AKLT ground states make good quantum wires.
 - ☞ Buffered measurements improve performance.
 - ☞ Simple projection operations make AKLTlike states even more so.
-
- ☞ Repeated gates: Markovian noise?
 - ☞ Two-qubit gates (CPHASE)? Thermal noise?
 - ☞ Analytic results (perturbative or variational using swap states)?
 - ☞ How does buffering relate to renormalization?

Thanks for your attention!